

Modular Arithmetic

Reduction modulo a polynomial $g(x)$ or modulo an integer m plays a central role in the mathematics of cryptography. Recall that for a monic polynomial $g(x)$ of positive degree, we define $a(x) \bmod g(x)$ to be the unique polynomial $a_0(x)$ with $\deg a_0(x) < \deg g(x)$ such that

$$a(x) = a_0(x) + a_1(x)g(x).$$

For an integer m , we define $a \bmod m$ to be the unique integer a_0 with $0 \leq a_0 < m$ such that $a = a_0 + a_1m$.

Fermat's little theorem. If p is a prime, then the relation $a^{p-1} \equiv 1 \pmod p$ holds for any integer a not divisible by p .

Note. The Magma function `mod` is the binary operator, with the syntax:

```
> m := 101;  
> 2^31 mod m;  
34
```

The same result can be achieved with the `Modexp`, or modular exponentiation function:

```
> Modexp(2,31,m);  
34
```

2. Let p be the prime $2^{31} - 1 = 2147483647$. Use the Magma function `Modexp` to verify Fermat's little theorem for several values of a . *Why would it be a bad idea to compute a^{p-1} and then reduce modulo p ?*

Chinese remainder theorem. Let p and q be distinct primes, then for every integer a and b there exists a unique integer c with $0 \leq c < pq$ such that $c \equiv a \pmod p$ and $c \equiv b \pmod q$.

If a , b , and c are as above, then for any integral polynomial $f(x)$, the integer $f(c)$ satisfies $f(c) \equiv f(a) \pmod p$ and $f(c) \equiv f(b) \pmod q$. Therefore $f(c) \bmod pq$ is the unique solution to the Chinese remainder theorem.

3. Let p be as above and let $q = (2^{61} + 1)/3 = 768614336404564651$. Compute $a^{p-1} \bmod pq$ for various primes using `Modexp`. Then reduce the result modulo p . How do you explain the result in terms of the Chinese remainder theorem and Fermat's little theorem?

Analogues of Fermat's little theorem also hold for polynomials.

Polynomial analogue of Fermat. If $g(x)$ is an irreducible polynomial of degree n over \mathbb{F}_2 , then the relation $a(x)^{2^n-1} \equiv 1 \pmod{g(x)}$ holds for any polynomial $a(x)$ not divisible by $g(x)$.

Chinese remainder theorem. Let $g(x)$ and $h(x)$ be monic polynomials with no common factors. Given any polynomials $a(x)$ and $b(x)$, there exists a unique polynomial $c(x)$ such that $c(x) \equiv a(x) \pmod{g(x)}$ and $c(x) \equiv b(x) \pmod{h(x)}$.

We can create and work with polynomials over \mathbb{F}_2 as demonstrated by the following Magma code.

```
> F2 := FiniteField(2);
> P2<x> := PolynomialRing(F2);
> f := x^17 + x^5 + 1;
> Factorization(f);
[
<x^17 + x^5 + 1, 1>
]
```

4. Let $g(x) = x^{17} + x^5 + 1$, and use the function `Modexp` to verify the polynomial analogue of Fermat's little theorem for the polynomials x , $x^2 + x + 1$, etc.

5. Let $h(x) = x^{17} + x^{15} + x^{10} + x^5 + 1$ and compute $a(x)^{2^{17}-1} \pmod{g(x)h(x)}$ for various $a(x)$. What is the result reduced modulo $g(x)$? Why does the same not hold true for $a(x)^{2^{17}-1} \pmod{g(x)h(x)}$, reduced modulo $h(x)$?